Nonlinear Optimization

Moritz Diehl

Systems Control and Optimization Laboratory, University of Freiburg, Germany

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(slides jointly developed with Armin Nurkanović, Titus Quah, Katrin Baumgärtner, Jim Rawlings)

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1 Basic definitions

- 2 Some classification of optimization problems
- 3 Optimality conditions
- 4 Nonlinear programming algorithms

Optimization Problem		
$\min_{w \in \mathbb{R}^n} \ F(w)$	(1a)	
s.t. $G(w) = 0$	(1b)	
$H(w) \ge 0$	(1c)	

Optimization Problem	Terminology
$\min_{w \in \mathbb{R}^n} F(w) $ (1a) s.t. $G(w) = 0$ (1b) $H(w) \ge 0$ (1c)	$w \in \mathbb{R}^{n} \text{ - vector of decision variables}$ $F : \mathbb{R}^{n} \to \mathbb{R} \text{ - objective function}$ $G : \mathbb{R}^{n} \to \mathbb{R}^{n_{G}} \text{ - equality constraints}$ $H : \mathbb{R}^{n} \to \mathbb{R}^{n_{H}} \text{ - inequality constraints}$



Optimization Problem	Terminology
$\min_{w \in \mathbb{R}^n} F(w) $ (1a) s.t. $G(w) = 0$ (1b) $H(w) \ge 0$ (1c)	 w ∈ ℝⁿ - vector of decision variables F : ℝⁿ → ℝ - objective function G : ℝⁿ → ℝ^{n_G} - equality constraints H : ℝⁿ → ℝ^{n_H} - inequality constraints

- only in a few special cases a closed form solution exists
- if F, G, H are nonlinear and smooth, we speak of a *nonlinear programming problem (NLP)*
- usually we need iterative algorithms to find an approximate solution
- ▶ in NMPC, the problem depends on parameters that change every sampling time



Definition

The feasible set of the optimization problem (1) is defined as $\Omega = \{w \in \mathbb{R}^n \mid G(w) = 0, H(w) \ge 0\}$. A point $w \in \Omega$ is is called a feasible point.



In the example, the feasible set is the intersection of the two grey areas (halfspace and circle)

Basic definitions: global and local minimizer

Definition (Global Minimizer)

Point $w^* \in \Omega$ is a global minimizer of the NLP (1) if for all $w \in \Omega$ it holds that $F(w) \ge F(w^*)$.

Definition (Local Minimizer)

Point $w^* \in \Omega$ is a **local minimizer** of the NLP (1) if there exists a ball $\mathcal{B}_{\epsilon}(w^*) = \{w | ||w - w^*|| \le \epsilon\}$ with $\epsilon > 0$, such that for all $w \in \mathcal{B}_{\epsilon}(w^*) \cap \Omega$ it holds that $F(w) \ge F(w^*)$

The value $F(w^*)$ at a local/global minimizer w^* is called local/global minimum, or minimum value.



Convex sets

a key concept in optimization





A set Ω is said to be convex if for any w_1, w_2 and any $\theta \in [0, 1]$ it holds $\theta w_1 + (1 - \theta) w_2 \in \Omega$ Figure inspired by Figure 2.2 in S. Boyd and L. Vandenberghe. Convex optimization. Cambridge university press, 2004.

Convex functions



• A function $F: \Omega \to \mathbb{R}$ is convex if for every $w_1, w_2 \in \Omega \subset \mathbb{R}^n$ and $\theta \in [0, 1]$ it holds that

$$F(\theta w_1 + (1-\theta)w_2) \le \theta F(w_1) + (1-\theta)F(w_2)$$

F is concave if and only if -F is convex
F is convex if and only if the epigraph

$$\mathrm{epi}F = \{(w,t) \in \mathbb{R}^{n_w+1} \mid w \in \Omega, F(w) \le t\}$$

is a convex set





A convex optimization problem

 $\min_{w} F(w)$ s.t. G(w) = 0 $H(w) \ge 0$ An optimization problem is convex if the objective function F is convex and the feasible set Ω is convex.

- For convex problems, every locally optimal solution is globally optimal
- First order conditions are necessary and sufficient
- "...in fact, the great watershed in optimization isn't between linearity and nonlinearity, but convexity and nonconvexity." R. T. Rockafellar, SIAM Review, 1993

Outline of the lecture



1 Basic definitions

2 Some classification of optimization problems

3 Optimality conditions

4 Nonlinear programming algorithms

Some classification of optimization problems

Optimization problems can be:

- unconstrained $(\Omega = \mathbb{R}^n)$ or constrained $(\Omega \subset \mathbb{R}^n)$
- convex or nonconvex
- linear or nonlinear
- differentiable or nonsmooth
- continuous or (mixed-)integer
- finite or infinite dimensional



Some classification of optimization problems

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"... the main fact, which should be known to any person dealing with optimization models, is that in general, optimization problems are unsolvable." Yurii Nesterov, Lectures on Convex Optimization, 2018. ("solvable" refers to finding a global minimizer)

Class 1: Linear Programming (LP)





- convex optimization problem
- 1947: simplex method by G. Dantzig
- a solution is always at a vertex of the feasible set (possibly a whole facet if nonunique)
- very mature and reliable

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Class 2: Quadratic Programming (QP)



Quadratic Program (QP)

$$\min_{w \in \mathbb{R}^n} \frac{1}{2} w^\top Q w + g^\top w$$

s.t. $Aw - b = 0$
 $Cw - d \ge 0$

- depending on Q, can be convex and nonconvex
- solved online in linear model predictive control
- many good solvers: Gurobi, OSQP, HPIPM, qpOASES, OOQP, DAQP...
- subsproblems in nonlinear optimization

Class 3: Nonlinear Programming (NLP)



Nonlinear Rrogram (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$
 $H(w) \ge 0$

- can be convex and nonconvex
- solved with iterative Newton-type algorithms
- solved in nonlinear model predictive control

Class 4: Mathematical Programming with Complementarity Constraints



 $\min_{w_0, w_1, w_2} F(w)$ s.t. G(w) = 0 $H(w) \ge 0$ $0 \le w_1 \perp w_2 \ge 0$ $w = [w_0^\top, w_1^\top, w_2^\top]^\top, w_1 \perp w_1 \Leftrightarrow w_1^\top w_2 = 0$



- more difficult than standard nonlinear programming
- feasible set is inherently nonsmooth and nonconvex
- powerful modeling concept
- requires specialized theory and algorithms







- inherently nonconvex feasible set
- due to combinatorial nature, NP-hard even for linear F, G, H
- branch and bound, branch and cut algorithms based on iterative solution of relaxed continuous problems

Optimal Control Problem (OCP)

$$\min_{x(\cdot),u(\cdot)} \int_{0}^{T} L_{c}(x(t), u(t)) dt + E(x(T))$$

s.t. $x(0) = \bar{x}_{0}$
 $\dot{x}(t) = f_{c}(x(t), u(t))$
 $0 \ge h(x(t), u(t)), \ t \in [0, T]$
 $0 \ge r(x(T))$

- decision variables $x(\cdot)$, $u(\cdot)$ in infinite dimensional function space
- infinitely many constraints ($t \in [0, T]$)
- Smooth ordinary differential equation (ODE) $\dot{x}(t) = f_c(x(t), u(t))$
- more generally, dynamic model can be based on
 - differential algebraic equations (DAE)
 - partial differential equations (PDE)
 - nonsmooth ODE
 - stochastic ODE
- OCP can be convex or nonconvex
- all or some components of u(t) may take integer values (mixed-integer OCP)



Continuous-time OCP

$$\min_{x(\cdot),u(\cdot)} \int_{0}^{T} L_{c}(x(t), u(t)) dt + E(x(T))$$

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Direct methods like direct collocation, multiple shooting first discretize, then optimize.



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Direct methods like direct collocation, multiple shooting first discretize, then optimize.

Discrete-time OCP (an NLP)

$$\min_{x,u} \sum_{k=0}^{N-1} \ell(x_k, u_k) + E(x_N)$$

s.t. $x_0 = \bar{x}_0$
 $x_{k+1} = f(x_k, u_k)$
 $0 \ge h(x_k, u_k), \ k = 0, \dots, N-1$
 $0 \ge r(x_N)$

Variables $x = (x_0, \ldots, x_N)$ and $u = (u_0, \ldots, u_{N-1})$ can be summarized in vector $w = (x, u) \in \mathbb{R}^n$.



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Discrete time NMPC Problem (an NLP)

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Variables $x = (x_0, \ldots, x_N)$ and $u = (u_0, \ldots, u_{N-1})$ can be summarized in vector $w = (x, u) \in \mathbb{R}^n$.

Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$
 $H(w) \ge 0$



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Algebraic characterization of unconstrained local minimizers



Consider the unconstrained problem: $\min_{w \in \mathbb{R}^n} F(w)$ First-Order Necessary Condition of Optimality (FONC) (in convex case also sufficient)

 $w^* \text{ local optimizer } \quad \Rightarrow \quad \nabla F(w^*) = 0, \ w^* \text{ stationary point}$

Second-Order Necessary Condition of Optimality (SONC)

 w^* local minimizer $\Rightarrow \nabla^2 F(w^*) \succeq 0$

Algebraic characterization of unconstrained local minimizers



Consider the unconstrained problem: $\min_{w \in \mathbb{R}^n} F(w)$

First-Order Necessary Condition of Optimality (FONC) (in convex case also sufficient)

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Second-Order Necessary Condition of Optimality (SONC)

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Second-Order Sufficient Conditions of Optimality (SOSC)

 $abla F(w^*) = 0$ and $abla^2 F(w^*) \succ 0 \quad \Rightarrow \quad x^*$ strict local minimizer

 $\nabla F(w^*) = 0$ and $\nabla^2 F(w^*) \prec 0 \implies x^*$ strict local maximizer

no conclusion can be drawn in the case $\nabla^2 F(w^*)$ is indefinite

Types of stationary points





a stationary point w with $\nabla F(w) = 0$ can be a minimizer, a maximizer, or a saddle point

Optimality conditions - unconstrained

- necessary conditions: find a candidate point (or to exclude points)
- sufficient conditions: verify optimality of a candidate point



Optimality conditions - unconstrained

- necessary conditions: find a candidate point (or to exclude points)
- sufficient conditions: verify optimality of a candidate point
- a minimizer must satisfy SONC, but does not have to satisfy SOSC



Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$

Lagrangian function $\mathcal{L}(w,\lambda) := F(w) - \lambda^{\top} G(w)$



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Definition (LICQ)

A point w satisfies Linear Independence Constraint Qualification (LICQ) if and only if $\nabla G(w) := \frac{\partial G}{\partial w}(w)^{\top}$ is full column rank



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First-Order Necessary Conditions (in convex case also sufficient)

Let F, G in C^1 . If w^* is a (local) minimizer, and w^* satisfies LICQ, then there is a unique vector λ such that:

$$\begin{split} \nabla_w \mathcal{L}(w^*,\lambda^*) &= \nabla F(w^*) - \nabla G(w^*)\lambda = 0 & \text{dual feasibility} \\ \nabla_\lambda \mathcal{L}(w^*,\lambda^*) &= G(w^*) = 0 & \text{primal feasibility} \end{split}$$

for equality constrained optimization



Primal Problem

$$p^* = \min_{w \in \mathbb{R}^n} F(w)$$
 s.t. $G(w) = 0$

with Lagrangian $\mathcal{L}(w, \lambda) := F(w) - \lambda^{\top} G(w).$

Lagrange dual function $\mathcal{Q}(\lambda) := \inf_{w \in \mathbb{R}^n} \mathcal{L}(w, \lambda)$

• $Q(\lambda)$ - concave in λ by construction • $Q(\lambda) \le p^*$ for all $\lambda \in \mathbb{R}^{n_G}$
for equality constrained optimization

Primal Problem

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Dual Problem

$$d^* = \max_{\lambda \in \mathbb{R}^{n_G}} \mathcal{Q}(\lambda)$$

- ▶ weak duality: $d^* \le p^*$, always holds
- strong duality: d* = p*, only holds for some problems (e.g. convex ones)

for equality constrained optimization

Primal Problem

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Q(λ) - concave in λ by construction
 Q(λ) ≤ p* for all λ ∈ ℝ^{n_G}



Dual Problem

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Wolfe Dual (in convex case)

$$d^* = \max_{w \in \mathbb{R}^n, \lambda \in \mathbb{R}^{n_G}} \mathcal{L}(w, \lambda)$$
s.t. $\nabla_w \mathcal{L}(w, \lambda) = 0$

(w constrained by lower level optimality)

The Karush-Kuhn-Tucker (KKT) conditions

Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$
 $H(w) \ge 0$

 $\mathcal{L}(w,\lambda) = F(w) - \lambda^{\top} G(w) - \mu^{\top} H(w)$



The Karush-Kuhn-Tucker (KKT) conditions

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Definition (LICQ)

A point \boldsymbol{w} satisfies LICQ if and only if

 $\left[\nabla G\left(w\right),\quad\nabla H_{\mathbb{A}}\left(w\right)\right]$

is full column rank

Active set $\mathbb{A} = \{i \mid H_i(w) = 0\}$

The Karush-Kuhn-Tucker (KKT) conditions

Nonlinear Program (NLP)

$$\label{eq:relation} \begin{split} \min_{w \in \mathbb{R}^n} \, F(w) \\ \text{s.t.} \; G(w) &= 0 \\ H(w) &\geq 0 \end{split}$$

Definition (LICQ)

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Active set $\mathbb{A} = \{i \mid H_i(w) = 0\}$

Theorem (KKT conditions - FONC for constrained optimization)

Let F, G, H be C^1 . If w^* is a (local) minimizer and satisfies LICQ, then there are unique vectors λ^* and μ^* such that (w^*, λ^*, μ^*) satisfies:

$$\begin{split} \nabla_{w}\mathcal{L}\left(w^{*},\,\mu^{*},\,\lambda^{*}\right) &= 0, \quad \mu^{*} \geq 0, \\ G\left(w^{*}\right) &= 0, \quad H\left(w^{*}\right) \geq 0 \\ \mu_{i}^{*}H_{i}(w^{*}) &= 0, \quad \forall i \\ \end{split} \label{eq:gamma-constraint} \begin{array}{l} \text{dual feasibility} \\ \text{primal feasibility} \\ \text{complementary slackness} \\ \end{array}$$

Complementarity Conditions



$$\label{eq:complementarity conditions} \begin{split} 0 \geq \mu \perp H(w) \geq 0 \text{ form an L-shaped,} \\ \text{nonsmooth manifold.} \end{split}$$

• $H_i(w^*) > 0$ then $\mu_i^* = 0$, and H_i is inactive





 $\begin{array}{l} \mbox{Complementarity conditions} \\ 0 \geq \mu \perp H(w) \geq 0 \mbox{ form an L-shaped,} \\ \mbox{nonsmooth manifold.} \end{array}$

- ▶ $H_i(w^*) > 0$ then $\mu_i^* = 0$, and H_i is inactive
- ▶ $\mu_i^* > 0$ and $H_i(w) = 0$ then $H_i(w)$ is strictly active





- ▶ $H_i(w^*) > 0$ then $\mu_i^* = 0$, and H_i is inactive
- ▶ $\mu_i^* > 0$ and $H_i(w) = 0$ then $H_i(w)$ is strictly active
- ▶ $\mu_i^* = 0$ and $H_i(w) = 0$ then then $H_i(w)$ is weakly active





- ▶ $H_i(w^*) > 0$ then $\mu_i^* = 0$, and H_i is inactive
- ▶ $\mu_i^* > 0$ and $H_i(w) = 0$ then $H_i(w)$ is strictly active
- ▶ $\mu_i^* = 0$ and $H_i(w) = 0$ then then $H_i(w)$ is weakly active
- We define the active set A as the set of indices i of the active constraints



Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint







Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint



$$\label{eq:expansion} \begin{split} \min_{w \in \mathbb{R}^2} \, F(w) \\ \text{s.t.} \ H(w) \geq 0 \\ \blacktriangleright \ -\nabla F \text{ is the gravity} \end{split}$$



Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint



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Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint



 $\min_{w \in \mathbb{R}^2} F(w)$ s.t. $H(w) \ge 0$

- \blacktriangleright $-\nabla F$ is the gravity
- ▶ $\mu \nabla H$ is the force of the fence. Sign $\mu \ge 0$ means the fence can only "push" the ball



Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint

 $\min_{w \in \mathbb{R}^2} F(w)$ s.t. H(w) > 0

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- $\blacktriangleright \nabla H$ gives the direction of the force and μ adjusts the magnitude



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- $\blacktriangleright \nabla H$ gives the direction of the force and μ adjusts the magnitude
- weakly active constraint: H (w) = 0, µ = 0 the ball touches the fence but no force is needed



$$\nabla \mathcal{L}(w,\mu) = \nabla F(w) - \mu \nabla H(w) = 0$$

Animation inspired by Lecture 2 of the Winter School on Numerical Optimal Control with Differential Algebraic Equations by S. Gros and M. Diehl, Freiburg, 2016.

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Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint

 $\min_{w \in \mathbb{R}^2} F(w)$
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- $\blacktriangleright \nabla H$ gives the direction of the force and μ adjusts the magnitude
- weakly active constraint: H (w) = 0, µ = 0 the ball touches the fence but no force is needed
- inactive constraint $H(w) > 0, \ \mu = 0$

$$H\left(w\right)>0,\quad \mu=0$$



$$\nabla \mathcal{L}(w,\mu) = \nabla F(w) - \mu \nabla H(w) = 0$$

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To solve a nonlinear system, solve a sequence of linear systems





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w

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w

To solve a nonlinear system, solve a sequence of linear systems





Nonlinear Optimization

w

In direct methods, we have to solve the discretized optimal control problem, which is a Nonlinear Program (NLP)

General Nonlinear Program (NLP)

$$\min_w F(w) \text{ s.t. } \begin{cases} G(w) = 0\\ H(w) \ge 0 \end{cases}$$

We first treat the case without inequalities

NLP only with equality constraints

$$\min_{w} F(w) \quad \text{s.t.} \quad G(w) = 0$$

Lagrange function

$$\mathcal{L}(w,\lambda) = F(w) - \lambda^T G(w)$$

Then for an optimal solution w^* exist multipliers λ^* such that

Nonlinear root-finding problem

$$\nabla_w \mathcal{L}(w^*, \lambda^*) = 0$$

$$G(w^*) = 0$$



Newton's method to solve

$$\nabla_w \mathcal{L}(w^*, \lambda^*) = 0$$

$$G(w^*) = 0 ?$$

results, at iterate $(\boldsymbol{w}^k,\boldsymbol{\lambda}^k),$ in the following linear system:

$$\begin{aligned} \nabla_w \mathcal{L}(w^k, \lambda^k) &+ \nabla^2_w \mathcal{L}(w^k, \lambda^k) \Delta w &- \nabla_w G(w^k) \Delta \lambda &= 0 \\ G(w^k) &+ \nabla_w G(w^k)^T \Delta w &= 0 \end{aligned}$$

Due to $\nabla \mathcal{L}(w^k,\lambda^k) = \nabla F(w^k) - \nabla G(w^k)\lambda^k$ this is equivalent to

$$\begin{aligned} \nabla_w F(w^k) &+ \nabla^2_w \mathcal{L}(w^k,\lambda^k) \Delta w &- \nabla_w G(w^k) \lambda^+ &= 0 \\ G(w^k) &+ \nabla_w G(w^k)^T \Delta w &= 0 \end{aligned}$$

with the shorthand $\lambda^+ = \lambda^k + \Delta \lambda$

Conditions

$$\begin{array}{rcl} \nabla_w F(w^k) & + \nabla^2_w \mathcal{L}(w^k, \lambda^k) \Delta w & - \nabla_w G(w^k) \lambda^+ &= 0 \\ G(w^k) & + \nabla_w G(w^k)^T \Delta w &= 0 \end{array}$$

are optimality conditions of a quadratic program (QP), namely:

Quadratic program

$$\begin{split} & \min_{\Delta w} \quad \nabla F(w^k)^T \Delta w + \frac{1}{2} \Delta w^T A^k \Delta w \\ & \text{s.t.} \qquad G(w^k) + \nabla G(w^k)^T \Delta w = 0, \end{split}$$

with $A^k = \nabla^2_w \mathcal{L}(w^k,\lambda^k)$



The full step Newton's Method iterates by solving in each iteration the Quadratic Progam

Quadratic Program in Sequential Quadratic Programming (SQP)

S

$$\min_{\Delta w} \quad \nabla F(w^k)^T \Delta w + \frac{1}{2} \Delta w^T A^k \Delta w$$

.t.
$$G(w^k) + \nabla G(w^k)^T \Delta w = 0,$$

with $A^k = \nabla^2_w \mathcal{L}(w^k, \lambda^k)$.

This obtains as solution the step Δw^k and the new multiplier $\lambda_{\rm QP}^+=\lambda^k+\Delta\lambda^k$

New iterate

This is the "full step, exact Hessian SQP method for equality constrained optimization".



Regard again NLP with both, equalities and inequalities:

NLP with equality and inequality constraints

$$\min_w F(w) \text{ s.t. } \begin{cases} G(w) = 0\\ H(w) \ge 0 \end{cases}$$

Lagrangian function for NLP with equality and inequality constraints

$$\mathcal{L}(w,\lambda,\mu) = F(w) - \lambda^T G(w) - \mu^T H(w)$$

Theorem (Karush-Kuhn-Tucker (KKT) conditions)

Let F, G, H be C^2 . If w^* is a (local) minimizer and satisfies LICQ, then there are unique vectors λ^* and μ^* such that (w^*, λ^*, μ^*) satisfies:

$$\nabla_w \mathcal{L} (w^*, \mu^*, \lambda^*) = 0$$

$$G (w^*) = 0$$

$$H(w^*) \ge 0$$

$$\mu^* \ge 0$$

$$H(w^*)^\top \mu^* = 0$$

► Last three "complementarity conditions" are nonsmooth

▶ Thus, this system cannot be solved by Newton's Method. But still with SQP...

By linearizing all functions and setting $\lambda^+ = \lambda^k + \Delta \lambda$, $\mu^+ = \mu^k + \Delta \mu$, we obtain the KKT conditions of the following Quadratic Program (QP)

Inequality Constrained Quadratic Program within SQP method

$$\begin{split} \min_{\Delta w} \quad \nabla F(w^k)^T \Delta w &+ \frac{1}{2} \Delta w^T A^k \Delta w \\ \text{s.t.} \quad \left\{ \begin{array}{ll} G(w^k) + \nabla G(w^k)^T \Delta w &= 0 \\ H(w^k) + \nabla H(w^k)^T \Delta w &\geq 0 \end{array} \right. \end{split}$$

with

$$A^k = \nabla^2_w \mathcal{L}(w^k, \lambda^k, \mu^k)$$

Its solution delivers the next SQP iterate

$$\Delta w^k, \quad \lambda_{\rm QP}^+, \quad \mu_{\rm QP}^+$$

Constrained Gauss-Newton Method

In special case of least squares objectives

Least squares objective function

$$F(w) = \frac{1}{2} \|R(w)\|_2^2$$

can approximate Hessian $\nabla^2_w \mathcal{L}(w^k,\lambda^k,\mu^k)$ by much cheaper

 $A^k = \nabla R(w) \nabla R(w)^\top.$

Need no multipliers to compute A^k .

Gauss-Newton QP = Constrained Linear Least Squares Problem

$$\begin{split} \min_{\Delta w} & \frac{1}{2} \| R(w^k) + \nabla R(w^k)^T \Delta w \|_2^2 \\ \text{s.t.} & G(w^k) + \nabla G(w^k)^T \Delta w = 0 \\ H(w^k) + \nabla H(w^k)^T \Delta w \ge 0 \end{split}$$

Linear convergence. Fast, if objective value $||R(w^*)||$ small or nonlinearity of R, G, H small



Interior Point Methods

(without equalities for simplicity of exposition)



NLP with inequalites

$$\min_{w} F(w)$$

s.t. $H(w) \ge 0$

KKT conditions

$$\nabla F(w) - \nabla H(w)^{\top} \mu = 0$$
$$0 \le \mu \perp H(w) \ge 0$$

Main difficulty: nonsmoothness of complementarity conditions



NLP with inequalites

$$\min_{w} F(w)$$

s.t. $H(w) \ge 0$

Idea: put inequality constraint into objective

NLP with inequalites

$$\min_{w} F(w)$$

s.t. $H(w) \ge 0$

Idea: put inequality constraint into objective

Barrier Problem

$$\min_{w} F(w) - \tau \sum_{i=1}^{m} \log(H_i(w)) \eqqcolon F_{\tau}(w)$$

Barrier Problem in Interior Point Method

NLP with inequalites

$$\min_{w} F(w)$$

s.t. $H(w) \ge 0$

Idea: put inequality constraint into objective

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Barrier Problem in Interior Point Method

NLP with inequalites

$$\min_{w} F(w)$$

s.t. $H(w) \ge 0$

Idea: put inequality constraint into objective

$$\min_{w} F(w) - \tau \sum_{i=1}^{m} \log(H_i(w)) \eqqcolon F_{\tau}(w)$$







Example NLP

$$\min_{w} 0.5w^2 - 2w$$

s.t. $-1 \le w \le 1$

$$\min_{w} \ 0.5w^2 - 2 - \tau \log(w+1) - \tau \log(1-w)$$





Example NLP

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Alternative interpretation

Barrier problem

$$\min_{w} F(w) - \tau \sum_{i=1}^{m} \log(H_i(w)) \eqqcolon F_{\tau}(w)$$

KKT conditions

$$\nabla F(w) - \tau \sum_{i=1}^{m} \frac{1}{H_i(w)} \nabla H_i(w) = 0$$

Introduce variable $\mu_i = \frac{\tau}{H_i(w)}$

$$\nabla F(w) - \nabla H(w)^{\top} \mu = 0$$
$$H_i(w)\mu_i = \tau$$
$$(H_i(w) > 0, \mu_i > 0)$$





Alternative interpretation

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Nonlinear programming problem

$$\min_{w,s} F(w)$$

s.t. $G(w) = 0$
 $H(w) - s = 0$
 $s \ge 0$

Smoothed KKT conditions

$$R_{\tau}(w, s, \lambda, \mu) = \begin{bmatrix} \nabla_{w} \mathcal{L}(w, \lambda, \mu) \\ G(w) \\ H(w) - s \\ \text{diag}(s)\mu - \tau e \end{bmatrix} = 0$$
$$(s, \mu > 0)$$

 $e = (1, \ldots, 1)$

Fix τ , perform Newton iterations

$$R_{\tau}(w, s, \lambda, \mu) + \nabla R_{\tau}(w, s, \lambda, \mu)^{\top} \Delta z = 0$$

with $z=(w,s,\lambda,\mu)$

Line-search

Find $\alpha \in (0,1)$

$$w^{k+1} = w^k + \alpha \Delta w$$
$$s^{k+1} = s^k + \alpha \Delta s$$
$$\lambda^{k+1} = \lambda^k + \alpha \Delta \lambda$$
$$\mu^{k+1} = \mu^k + \alpha \Delta \mu$$

such that $s^{k+1}>0, \mu^{k+1}>0$

Reduce $\boldsymbol{\tau},$ and perform next Newton iterations solve, etc





- ▶ optimization problem come in many variants (LP, QP, NLP, MPCC, MINLP, OCP,)
- each problem class be addressed with suitable software
- nonlinear MPC needs to solve nonlinear programs (NLP)
- Lagrangian function, duality, and KKT conditions are important concepts
- ▶ for convex problems holds strong duality, KKT conditions sufficient for global optimality
- Newton-type optimization for NLP solves the nonsmooth KKT conditions via Sequential Quadratic Programming (SQP, e.g. acados) or via Interior Point Method (e.g. ipopt)
- ▶ NLP solvers need to evaluate first and second order derivatives (e.g. via CasADi)

Some References



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