1. Theory and algorithms for nonlinear programming

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1 Basic definitions

- 2 Some classifications of optimization problems
- 3 Optimality conditions
- 4 Nonlinear programming algorithms

What is an optimization problem?

Optimization is a powerful tool used in all quantitative sciences.



Minimize (or maximize) an objective function F(w) depending on decision variables w subject to equality and/or inequality constrains

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An optimization problem	Terminology
$\min_{w \in \mathbb{R}^n} F(w) $ (1a) s.t. $G(w) = 0$ (1b) $H(w) \ge 0$ (1c)	 <i>w</i> ∈ ℝⁿ - decision variable <i>F</i> : ℝⁿ → ℝ - objective <i>G</i> : ℝⁿ → ℝ^{n_G} - equality constraints <i>H</i> : ℝⁿ → ℝ^{n_H} - inequality constraints

- ▶ If F, G, H are nonlinear and smooth, we speak of a *nonlinear programming problem* (NLP).
- Only in few special cases a closed form solution exists.
- Use an iterative algorithm to find an approximate solution.
- Problem may be parametric, and some (or all) functions depend on a fixed parameter $p \in \mathbb{R}^p$, e.g. model predictive control.



Definition

The feasible set of the optimization problem (1) is defined as $\Omega = \{w \in \mathbb{R}^n \mid G(w) = 0, H(w) \ge 0\}$. A point $w \in \Omega$ is is called a feasible point.



In the example, the feasible set is the intersection of the two grey areas (halfspace and circle).

Basic definitions: local and global minimizer

Definition (Local minimizer)

A point $w^* \in \Omega$ is called a **local minimizer** of the optimization problem (1) if there exists an open ball $\mathcal{B}_{\epsilon}(w^*)$ with $\epsilon > 0$, such that for all $w \in \mathcal{B}_{\epsilon}(w^*) \cap \Omega$ it holds that $F(w) \ge F(w^*)$.

Definition (Global minimizer)

A point $w^* \in \Omega$ is called a **global minimizer** of (1) if for all $w \in \Omega$ it holds that $F(w) \ge F(w^*)$.

The value F(w*) at a local/global minimizer w* is called local/global minimum.



Convex sets

A key concept in optimization is convexity.



A set Ω is said to be convex if for any w_1, w_2 and any $\theta \in [0, 1]$ it holds $\theta w_1 + (1 - \theta) w_2 \in \Omega$



Figure inspired by Figure 2.2 in S. Boyd and L. Vandenberghe. Convex optimization. Cambridge university press, 2004.

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Convex functions



A function F is convex if for every $w_1, w_2 \in \mathbb{R}^n$ and $\theta \in [0, 1]$ it holds that

 $F(\theta w_1 + (1-\theta)w_2) \le \theta F(w_1) + (1-\theta)F(w_2)$

F is concave if and only if -F is convex
F is convex if and only if the epigraph

$$epiF = \{(w,t) \in \mathbb{R}^{n_w+1} \mid F(w) \le t\}$$

is a convex set



w

A convex optimization problem

 $\min_{w \in \mathbb{R}^n} F(w)$ s.t. G(w) = 0 $H(w) \ge 0$ An optimization problem is convex if the objective function F is convex and the feasible set Ω is convex.

- For convex problems, every locally optimal solution is globally optimal.
- First-order optimality conditions are necessary and sufficient.
- Many iterative algorithms for nonconvex optimization solve a sequence of convex optimization problems.

A convex optimization problem

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- For convex problems, every locally optimal solution is globally optimal.
- First-order optimality conditions are necessary and sufficient.
- Many iterative algorithms for nonconvex optimization solve a sequence of convex optimization problems.
- "...in fact, the great watershed in optimization isn't between linearity and nonlinearity, but convexity and nonconvexity." R. T. Rockafellar, SIAM Review, 1993

Outline of the lecture



1 Basic definitions

2 Some classifications of optimization problems

3 Optimality conditions

4 Nonlinear programming algorithms

Some classifications of optimization problems

Optimization problems can be:

- unconstrained $(\Omega = \mathbb{R}^n)$ or constrained $(\Omega \subset \mathbb{R}^n)$
- convex or nonconvex
- linear or nonlinear
- differentiable or nonsmooth
- continuous or (mixed-)integer
- finite or infinite dimensional



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"... the main fact, which should be known to any person dealing with optimization models, is that in general, optimization problems are unsolvable." Yurii Nesterov, Lectures on Convex Optimization, 2018. ("solvable" refers to finding a global minimizer)

Class 1: Linear Programming (LP)





- convex optimization problem
- 1947: simplex method by Dantzig, 1984: polynomial time interior-point method by Karmarkar
- a solution is always at a vertex of the feasible set (possibly a whole facet if nonunique)
- very mature and reliable

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Class 2: Quadratic Programming (QP)





- depending on Q, can be convex and nonconvex
- many good solvers: OSQP, HPIPM, qpOASES, Gurobi, Clarabel, DAQP, OOQP, MOSEK, ...
- solved online in linear model predictive control
- subsproblems in nonlinear optimization

Class 3: Nonlinear Program (NLP)





- \blacktriangleright F, G, H smooth functions, can be convex and nonconvex
- solved with iterative Newton-type algorithms
- solved in nonlinear model predictive control

 w_1

Class 4: Mathematical programs with Complementarity Constraints $_{\mbox{(MPCC)}}$





- more difficult than standard nonlinear programming
- feasible set is inherently nonsmooth and nonconvex
- powerful modeling concept
- requires specialized theory and algorithms (Lectures 5 and 6 on Wednesday)

Class 5: Mixed-integer programming







- inherently nonconvex feasible set
- due to combinatorial nature, NP-hard even for linear F, G, H
- branch and bound, branch and cut algorithms based on iterative solution of relaxed continuous problems

Continuous-time Optimal Control Problem

$$\begin{split} \min_{x(\cdot),u(\cdot)} & \int_{0}^{T} L_{\rm c}(x(t),u(t)) \, \mathrm{d}t + E(x(T)) \\ \text{s.t.} & x(0) = \bar{x}_{0} \\ & \dot{x}(t) = f_{\rm c}(x(t),u(t)) \\ & 0 \geq h(x(t),u(t)), \ t \in [0,T] \\ & 0 \geq r(x(T)) \end{split}$$

- \blacktriangleright decision variables $x(\cdot), \, u(\cdot)$ in infinite dimensional function space
- infinitely many constraints for $t \in [0,T]$
- smooth ordinary differential equations (ODE)

 $\dot{x}(t) = f_{\rm c}(x(t), u(t))$

- more generally, dynamic model can be based on
 - differential algebraic equations (DAE)
 - partial differential equations (PDE)
 - stochastic ODE
 - nonsmooth ODE (treated on Tuesday)
- OCP can be convex or nonconvex
- all or some components of u(t) may take integer values (mixed-integer OCP)

Direct optimal control methods solve Nonlinear Programs (NLP)

Treated in detail in the 2nd lecture.

Continuous time OCP

$$\min_{x(\cdot),u(\cdot)} \int_{0}^{T} L_{c}(x(t), u(t)) dt + E(x(T))$$

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Direct methods like direct collocation, multiple shooting first discretize, then optimize.

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Direct methods like direct collocation, multiple shooting first discretize, then optimize.

Discrete time OCP (an NLP)

$$\min_{x,u} \sum_{k=0}^{N-1} \ell(x_k, u_k) + E(x_N)$$

s.t. $x_0 = \bar{x}_0$
 $x_{k+1} = f(x_k, u_k)$
 $0 \ge h(x_k, u_k), \ k = 0, \dots, N-1$
 $0 \ge r(x_N)$

Variables $x = (x_0, \ldots, x_N)$ and $u = (u_0, \ldots, u_{N-1})$ can be summarized in vector $w = (x, u) \in \mathbb{R}^n$.

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Discrete time NMPC Problem (an NLP)

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Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$
 $H(w) \ge 0$



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Algebraic characterization of **unconstrained** local optima





Second-Order **Necessary** Condition of Optimality (SONC)

 $w^* \text{ local minimizer } \quad \Rightarrow \quad \nabla^2 F(w^*) \succeq 0$

Algebraic characterization of unconstrained local optima



Consider the unconstrained problem: $\min_{w \in \mathbb{R}^n} F(w)$ First-Order Necessary Condition of Optimality (FONC) (in convex case also sufficient) w^* local optimizer $\Rightarrow \nabla F(w^*) = 0, w^*$ stationary point

Second-Order Necessary Condition of Optimality (SONC)

 w^* local minimizer $\Rightarrow \nabla^2 F(w^*) \succeq 0$

Second-Order Sufficient Conditions of Optimality (SOSC)

 $abla F(w^*) = 0$ and $abla^2 F(w^*) \succ 0 \quad \Rightarrow \quad x^*$ strict local minimizer

 $\nabla F(w^*) = 0$ and $\nabla^2 F(w^*) \prec 0 \implies x^*$ strict local maximizer

no conclusion can be drawn in the case $\nabla^2 F(w^*)$ is indefinite.

Type of stationary points





a stationary point w^* with $\nabla F(w^*) = 0$ can be a minimizer, a maximizer, or a saddle point

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Optimality conditions - unconstrained

- Necessary conditions: find a candidate point (or to exclude points)
- Sufficient conditions: verify optimality of a candidate point



Optimality conditions - unconstrained

- Necessary conditions: find a candidate point (or to exclude points)
- Sufficient conditions: verify optimality of a candidate point
- A minimizer must satisfy SONC, but does not have to satisfy SOSC



Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$

Lagrangian function: $\mathcal{L}(w, \lambda) = F(w) - \lambda^{\top} G(w)$



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Definition (LICQ)

A point w satisfies Linear Independence Constraint Qualification (LICQ) if $\nabla G(w) \coloneqq \frac{\partial G(w)}{\partial w}^{\top}$ is full column rank.



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First-order necessary conditions (in convex case also sufficient)

Let F, G in C^1 . If w^* is a (local) minimizer, <u>and</u> w^* satisfies LICQ, then there is a unique vector λ such that:

$$\begin{split} \nabla_w \mathcal{L}(w^*,\lambda^*) &= \nabla F(w^*) - \nabla G(w^*)\lambda = 0 & \text{dual feasibility} \\ \nabla_\lambda \mathcal{L}(w^*,\lambda^*) &= G(w^*) = 0 & \text{primal feasibility} \end{split}$$

The Karush-Kuhn-Tucker (KKT) conditions

Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $G(w) = 0$
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Definition (LICQ)

A point \boldsymbol{w} satisfies LICQ if

 $\left[\nabla G\left(w\right), \quad \nabla H_{\mathcal{A}}\left(w\right)\right]$

is full column rank.

Active set $\mathcal{A} = \{i \mid H_i(w) = 0\}$
The Karush-Kuhn-Tucker (KKT) conditions

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Theorem (KKT conditions - FONC for constrained optimization)

Let F, G, H be C^1 . If w^* is a (local) minimizer and satisfies LICQ, then there are unique vectors λ^* and μ^* such that (w^*, λ^*, μ^*) satisfies:

$$\begin{split} \nabla_{w}\mathcal{L}\left(w^{*},\,\mu^{*},\,\lambda^{*}\right) &= 0, \quad \mu^{*} \geq 0, \\ G\left(w^{*}\right) &= 0, \quad H\left(w^{*}\right) \geq 0 \\ \mu_{i}^{*}H_{i}(w^{*}) &= 0, \quad \forall i \\ \end{split} \label{eq:gamma-constraint} \begin{array}{l} \text{dual feasibility} \\ \text{primal feasibility} \\ \text{complementary slackness} \\ \end{array}$$

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- complementarity conditions $0 \le \mu^* \perp H(w^*) \ge 0$
- ▶ same as $\min(\mu^*, H(w^*)) = 0$
- \blacktriangleright zero level set of min is an L-shaped set



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Cases:

▶ $H_i(w^*) > 0$ then $\mu_i^* = 0$, and $H_i(w)$ is inactive



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Cases:

- ▶ $H_i(w^*) > 0$ then $\mu_i^* = 0$, and $H_i(w)$ is inactive
- ▶ $\mu_i^* > 0$ and $H_i(w) = 0$ then $H_i(w)$ is strictly active



- complementarity conditions $0 \le \mu^* \perp H(w^*) \ge 0$
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Cases:

- ▶ $H_i(w^*) > 0$ then $\mu_i^* = 0$, and $H_i(w)$ is inactive
- ▶ $\mu_i^* > 0$ and $H_i(w) = 0$ then $H_i(w)$ is strictly active
- ▶ $\mu_i^* = 0$ and $H_i(w) = 0$ then then $H_i(w)$ is weakly active







Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint







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Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint

 $\min_{w \in \mathbb{R}^n} F(w)$ s.t. $H(w) \ge 0$

- \blacktriangleright $-\nabla F$ is the gravity
- ▶ $\mu \nabla H$ is the force of the fence. Sign $\mu \ge 0$ means the fence can only "push" the ball



Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint

 $\min_{w \in \mathbb{R}^n} F(w)$ s.t. H(w) > 0

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- ► ∇H gives the direction of the force and µ adjusts the magnitude.





Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint

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- ▶ active constraint: $H(w) = 0, \ \mu > 0$



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- ► ∇H gives the direction of the force and µ adjusts the magnitude.
- ▶ active constraint: $H(w) = 0, \ \mu > 0$
- weakly active constraint: $H(w) = 0, \ \mu = 0$ the ball touches the fence but no force is needed





$$\nabla \mathcal{L}(w,\mu) = \nabla F(w) - \nabla H(w)\mu = 0$$



Ball rolling down a valley blocked by a fence - test problem with two variables and one inequality constraint

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- ► ∇H gives the direction of the force and µ adjusts the magnitude.
- ▶ active constraint: $H(w) = 0, \ \mu > 0$
- weakly active constraint: $H(w) = 0, \ \mu = 0$ the ball touches the fence but no force is needed
- ▶ inactive constraint: $H(w) > 0, \ \mu = 0$



$$\nabla \mathcal{L}(w,\mu) = \nabla F(w) - \nabla H(w)\mu = 0$$



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To solve a nonlinear system, solve a sequence of linear systems





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Linearization of F at linearization point \bar{w}



To solve a nonlinear system, solve a sequence of linear systems



Linearization of F at linearization point \bar{w}



To solve a nonlinear system, solve a sequence of linear systems



Linearization of F at linearization point \bar{w}



In direct methods, we have to solve the discretized optimal control problem, which is a Nonlinear Program (NLP)

General Nonlinear Program (NLP)

$$\min_{w \in \mathbb{R}^n} F(w) \text{ s.t. } \begin{cases} G(w) = 0\\ H(w) \ge 0 \end{cases}$$

We first treat the case without inequality constraints

NLP only with equality constraints

$$\min_{w \in \mathbb{R}^n} F(w) \text{ s.t. } G(w) = 0$$



Lagrange function

$$\mathcal{L}(w,\lambda) = F(w) - \lambda^{\top} G(w)$$

For an optimal solution w^* there exist multipliers λ^* such that

Nonlinear root-finding problem

$$\nabla_w \mathcal{L}(w^*, \lambda^*) = 0$$

$$G(w^*) = 0$$

Use Newton's method to solve:

$$\nabla_w \mathcal{L}(w^*, \lambda^*) = 0$$

$$G(w^*) = 0 ?$$

Given an iterate (w^k, λ^k) , the linearization reads as:

$$\begin{aligned} \nabla_{w} \mathcal{L}(w^{k},\lambda^{k}) &+ \nabla_{w}^{2} \mathcal{L}(w^{k},\lambda^{k}) \Delta w &- \nabla_{w} G(w^{k}) \Delta \lambda &= 0 \\ G(w^{k}) &+ \nabla_{w} G(w^{k})^{\top} \Delta w &= 0 \end{aligned}$$

Due to $\nabla \mathcal{L}(w^k,\lambda^k) = \nabla F(w^k) - \nabla G(w^k)\lambda^k$, this is equivalent to:

$$\begin{aligned} \nabla_w F(w^k) &+ \nabla^2_w \mathcal{L}(w^k, \lambda^k) \Delta w &- \nabla_w G(w^k) \lambda^+ &= 0 \\ G(w^k) &+ \nabla_w G(w^k)^\top \Delta w &= 0 \end{aligned}$$

with the shorthand $\lambda^+ = \lambda^k + \Delta \lambda$

Conditions

$$\begin{array}{rcl} \nabla_w F(w^k) & + \nabla^2_w \mathcal{L}(w^k, \lambda^k) \Delta w & - \nabla_w G(w^k) \lambda^+ &= 0 \\ G(w^k) & + \nabla_w G(w^k)^\top \Delta w &= 0 \end{array}$$

are the KKT optimality conditions of a quadratic program (QP), namely:

Quadratic program

$$\min_{\Delta w \in \mathbb{R}^n} \quad \nabla F(w^k)^\top \Delta w + \frac{1}{2} \Delta w^\top A^k \Delta w$$

s.t.
$$G(w^k) + \nabla G(w^k)^\top \Delta w = 0,$$

with $A^k = \nabla^2_w \mathcal{L}(w^k, \lambda^k)$





The full step Newton's Method iterates by solving in each iteration the QP

Quadratic program in Sequential Quadratic Programming (SQP)

$$\min_{\Delta w \in \mathbb{R}^n} \quad \nabla F(w^k)^\top \Delta w + \frac{1}{2} \Delta w^\top A^k \Delta w$$

s.t.
$$G(w^k) + \nabla G(w^k)^\top \Delta w = 0,$$

with $A^k = \nabla^2_w \mathcal{L}(w^k,\lambda^k)$

This obtains as solution the step Δw^k and the new multiplier $\lambda_{\rm QP}^+ = \lambda^k + \Delta \lambda^k$

New iterate

This is the "full step, exact Hessian SQP method for equality constrained optimization".



Regard again NLP with both, equality and inequality constraints:

NLP with equality and inequality constraints

$$\min_{w \in \mathbb{R}^n} F(w) \text{ s.t. } \begin{cases} G(w) = 0\\ H(w) \ge 0 \end{cases}$$

Lagrangian function for NLP with equality and inequality constraints

$$\mathcal{L}(w,\lambda,\mu) = F(w) - \lambda^{\top} G(w) - \mu^{\top} H(w)$$

Theorem (Karush-Kuhn-Tucker (KKT) conditions)

Let F, G, H be C^2 . If w^* is a (local) minimizer and satisfies LICQ, then there are unique vectors λ^* and μ^* such that (w^*, λ^*, μ^*) satisfies:

$$\nabla_w \mathcal{L} (w^*, \mu^*, \lambda^*) = 0$$

$$G (w^*) = 0$$

$$H(w^*) \ge 0$$

$$\mu^* \ge 0$$

$$H(w^*)^\top \mu^* = 0$$

► Last tree *complementarity conditions* make the KKT conditions nonsmooth

▶ This system cannot be solved by plain Newton's method. But we can use SQP...

By Linearizing all functions within the KKT Conditions, and setting $\lambda^+ = \lambda^k + \Delta \lambda$ and $\mu^+ = \mu^k + \Delta \mu$, we obtain the KKT conditions of a Quadratic Program (QP)

QP with inequality constraints

$$\min_{w \in \mathbb{R}^n} \quad \nabla F(w^k)^\top \Delta w + \frac{1}{2} \Delta w^\top A^k \Delta w$$

s.t.
$$G(w^k) + \nabla G(w^k)^\top \Delta w = 0$$

$$H(w^k) + \nabla H(w^k)^\top \Delta w \ge 0$$

with $A^k = \nabla^2_w \mathcal{L}(w^k, \lambda^k, \mu^k)$

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QP with inequality constraints

$$\min_{w \in \mathbb{R}^n} \quad \nabla F(w^k)^\top \Delta w + \frac{1}{2} \Delta w^\top A^k \Delta w$$

s.t.
$$G(w^k) + \nabla G(w^k)^\top \Delta w = 0$$

$$H(w^k) + \nabla H(w^k)^\top \Delta w \ge 0$$

with $A^k = \nabla^2_w \mathcal{L}(w^k, \lambda^k, \mu^k)$

► QP solution: Δw^k , λ_{QP}^+ , μ_{QP}^+ ► full step: $w^{k+1} = w^k + \Delta w^k$, $\lambda^{k+1} = \lambda_{\text{QP}}^+$, $\mu^{k+1} = \mu_{\text{QP}}^+$

nonsmooth complementarity conditions resolved at QP level

Interior-point methods

(without equality constraint for lighter notation)



NLP with inequalites

$$\min_{w \in \mathbb{R}^n} \quad F(w)$$
 s.t.
$$H(w) \ge 0$$

KKT conditions

 $\nabla F(w) - \nabla H(w)\mu = 0$ $0 \le \mu \perp H(w) \ge 0$

- Main difficulty: nonsmoothness of complementarity conditions
- 4th lecture (Tuesday) will show why Newton's method does not work for nonsmooth problems



NLP with inequalites

$$\min_{v \in \mathbb{R}^n} \quad F(w)$$

s.t. $H(w) \ge 0$

Idea: put inequality constraint into objective

$$\min_{w \in \mathbb{R}^n} F(w) - \tau \sum_{i=1}^m \log(H_i(w)) \rightleftharpoons F_\tau(w)$$



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$$\chi(H_i(w)) = \begin{cases} 0 & \text{ if } H_i(w) \ge 0\\ \infty & \text{ if } H_i(w) < 0 \end{cases}$$

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 $\ensuremath{\mathsf{NLP}}$ with inequalites

11

$$\min_{w \in \mathbb{R}^n} F(w)$$

s.t. $H(w) \ge 0$

Idea: put inequality constraint into objective

$$\min_{w \in \mathbb{R}^n} F(w) - \tau \sum_{i=1}^m \log(H_i(w)) \rightleftharpoons F_\tau(w)$$



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Example NLP

$$\min_{v \in \mathbb{R}^2} \quad 0.5w^2 - 2w$$

s.t.
$$-1 \le w \le 1$$

$$\min_{w \in \mathbb{R}^2} \ 0.5w^2 - 2 - \tau \log(w+1) - \tau \log(1-w)$$





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Alternative interpretation

Barrier problem

$$\min_{w \in \mathbb{R}^n} F(w) - \tau \sum_{i=1}^m \log(H_i(w)) \eqqcolon F_\tau(w)$$

KKT conditions ($\nabla F_{\tau}(w) = 0$)

$$\nabla F(w) - \tau \sum_{i=1}^{m} \frac{1}{H_i(w)} \nabla H_i(w) = 0$$

Introduce variable $\mu_i = \frac{\tau}{H_i(w)}$

$$\nabla F(w) - \nabla H(w)^{\top} \mu = 0$$
$$H_i(w)\mu_i = \tau$$
$$(H_i(w) > 0, \mu_i > 0)$$





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Nonlinear programming problem

$$\min_{\substack{w \in \mathbb{R}^n \\ \in \mathbb{R}^{n_H}}} F(w)$$

s.t. $G(w) = 0$
 $H(w) - s = 0$
 $s \ge 0$

Smoothed KKT conditions

$$R_{\tau}(w, s, \lambda, \mu) = \begin{bmatrix} \nabla_{w} \mathcal{L}(w, \lambda, \mu) \\ G(w) \\ H(w) - s \\ \operatorname{diag}(s)\mu - \tau e \end{bmatrix} = 0$$
$$(s, \mu > 0)$$

Solve approximately with Newton's method for fixed $\boldsymbol{\tau}$

$$R_{\tau}(w, s, \lambda, \mu) + \nabla R_{\tau}(w, s, \lambda, \mu) \Delta z = 0$$

with $z=(w,s,\lambda,\mu)$

Line-serach

Find $\alpha \in (0,1)$

$$w^{k+1} = w^k + \alpha \Delta w$$
$$s^{k+1} = s^k + \alpha \Delta s$$
$$\lambda^{k+1} = \lambda^k + \alpha \Delta \lambda$$
$$\mu^{k+1} = \mu^k + \alpha \Delta \mu$$

such that $s^{k+1}>0, \mu^{k+1}>0$

and reduce τ ...

 $e = (1, \ldots, 1)$

A. Nurkanović





- Optimization problem come in many variants (LP, QP, NLP, MPCC, MINLP, OCP,)
- Each problem class be addressed with suitable software.
- Lagrangian function, duality, and KKT conditions are important concepts
- ► For convex problems KKT conditions sufficient for global optimality.
- Newton-type optimization for NLP solves the nonsmooth KKT conditions via Sequential Quadratic Programming (SQP) or via the Interior-Point Method.
- NLP solvers need to evaluate first and second order derivatives (e.g. via CasADi).

- Duality for convex optimization problems.
- First-order methods (gradient descent, stochastic gradient descent, ...).
- Solution methods for linear and quadratic programs (active set, interior-point, simplex).
- Augmented Lagrangian methods for constrained optimization.
- Solution methods for mixed-integer problems (branch and bound, ...)
- Computing derivatives via automatic differentiation.
- Globalization strategies (linear search vs trust region, merit functions vs filter).
- Regularization (convexification of the Hessian, LICQ violation).



Nonlinear optimization:

- Nocedal, Jorge, and Stephen J. Wright, eds. Numerical optimization. New York, NY: Springer New York, 2006.
- Biegler, Lorenz T. Nonlinear programming: concepts, algorithms, and applications to chemical processes. Society for Industrial and Applied Mathematics, 2010.

Convex optimization:

- Boyd, Stephen, and Lieven Vandenberghe. Convex Optimization. Cambridge University Press, 2004. online: https://web.stanford.edu/~boyd/cvxbook/
- Rockafellar, R. T., Fundamentals of optimization. Lecture Notes 2007. online: https://sites.math.washington.edu/~rtr/fundamentals.pdf



Numerical optimization video lectures by Moritz Diehl (highly recommended!):

- Videos: https://www.syscop.de/teaching/ws2020/numerical-optimization
- Lecture notes: https://publications.syscop.de/Diehl2016.pdf

Lecture notes/slides by Mario Zanon Sébastien Gros

https://mariozanon.wordpress.com/teaching/ numerical-methods-for-optimal-control/

Optimization software:

- https://plato.asu.edu/guide.html
- https://www.syscop.de/research/software



- Moritz Diehl, Sébastien Gros. "Numerical optimal control (Draft)," Lecture notes, 2024. online: https://www.syscop.de/files/2024ws/NOC/book-NOCSE.pdf
- Karmarkar, Narendra. "A new polynomial-time algorithm for linear programming." In Proceedings of the sixteenth annual ACM symposium on Theory of computing, pp. 302-311. 1984.
- Dantzig, George B. "Origins of the simplex method." In A history of scientific computing, pp. 141-151. 1990.

Optimality conditions for NLP with equality and/or inequality constraints:

- First-Order Necessary Conditions: A local optimizer of a (differentiable) NLP is a KKT point
- Second-Order Sufficient Conditions require positivity of the Hessian in all critical feasible directions

Nonconvex problem \Rightarrow a minimizer is not necessarily a global minimizer. Note: some nonconvex problems may have a unique **minimum**

Some important practical consequences...

- A KKT point may not be a local (global) optimizer ... the lack of equivalence results from a lack of regularity and/or SOSC
- A local (global) optimizer **may not** be a KKT point
 - ... due to violation of constraint qualifications, e.g. LICQ violated